

we don't have to decide
fairness ourselves

[<draft-briscoe-tsvwg-relax-fairness-00.txt>](#)

intent: build consensus then Informational

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IETF-70 tsvarea Dec 2007



shifting IETF focus from fairness to accountability

| | design-time | run-time |
|-----------------------------------|--|---|
| problem | IETF doesn't, can't and shouldn't decide fairness | |
| solution process | IETF's role: enable accountability for congestion | users, apps & operators can (optionally) make principled fairness choices |
| | IETF/IRTF can truly meet dynamic app req's and minimise congestion | |
| best metric: congestion volume | | |

this talk primarily about the technical problem

- fairness is run-time, IETF is design-time

fair bottleneck bit-rate?

two incompatible partial worldviews

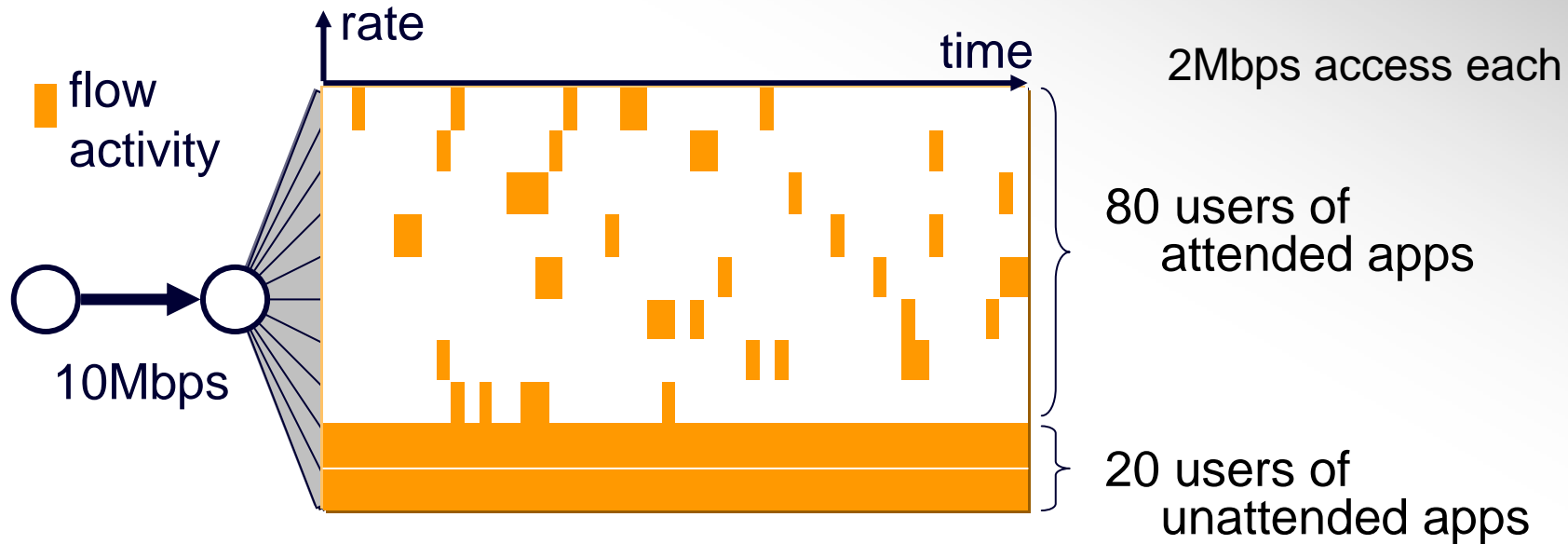
| 'flow rate equality' | 'volume accounting' |
|----------------------|---------------------|
| per flow | per user |
| instantaneous | over time |

- IETF aware that fairness should be per user
 - per flow is reasonable approx'n if users open similar no's of flows

base example

different activity factors

volume over peak period [B/hr]
 instantaneous equivalent: traffic intensity [b/s]
 = (bit rate when active [b/s]) x (activity factor [%])



| usage type | no. of users | activity factor | ave. simul flows /user | TCP bit rate /user | vol/day (16hr) /user | traffic intensity /user |
|------------|--------------|-----------------|------------------------|--------------------|----------------------|-------------------------|
| attended | 80 | 10% | = | 357kbps | 257MB | 35.7kbps |
| unattended | 20 | 100% | = | 357kbps | 2570MB | 357kbps |

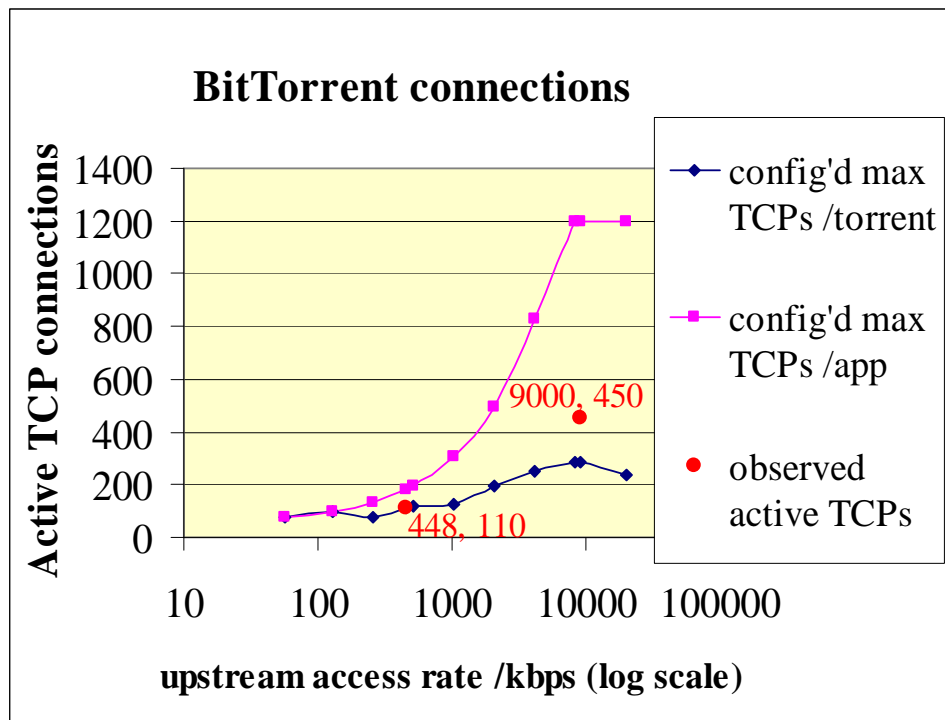
x1

x10

x10

realistic numbers?

there are elephants in the room



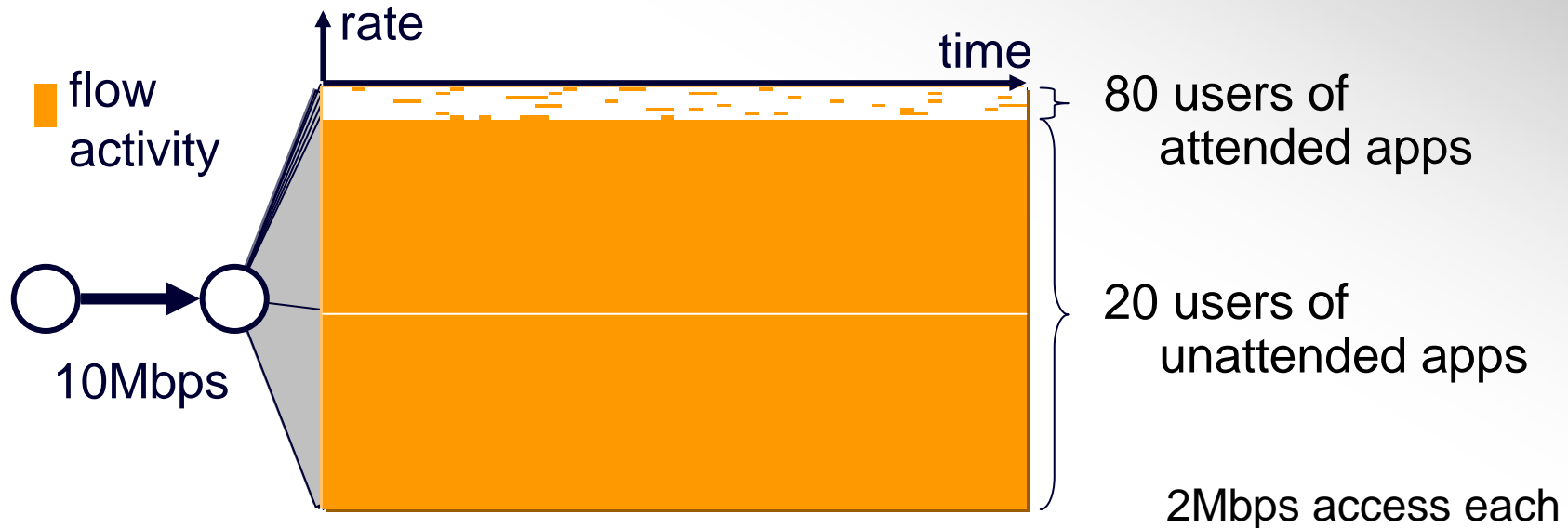
- number of TCP connections
 - Web1.1: **2**
 - BitTorrent: **~100**; see graph

details suppressed:

- users on spectrum of mixes of the two types
- utilisation never 100%
 - but near enough during peak period
- on DSL, upstream constrains most p2p apps
 - other access (fixed & wireless) more symmetric

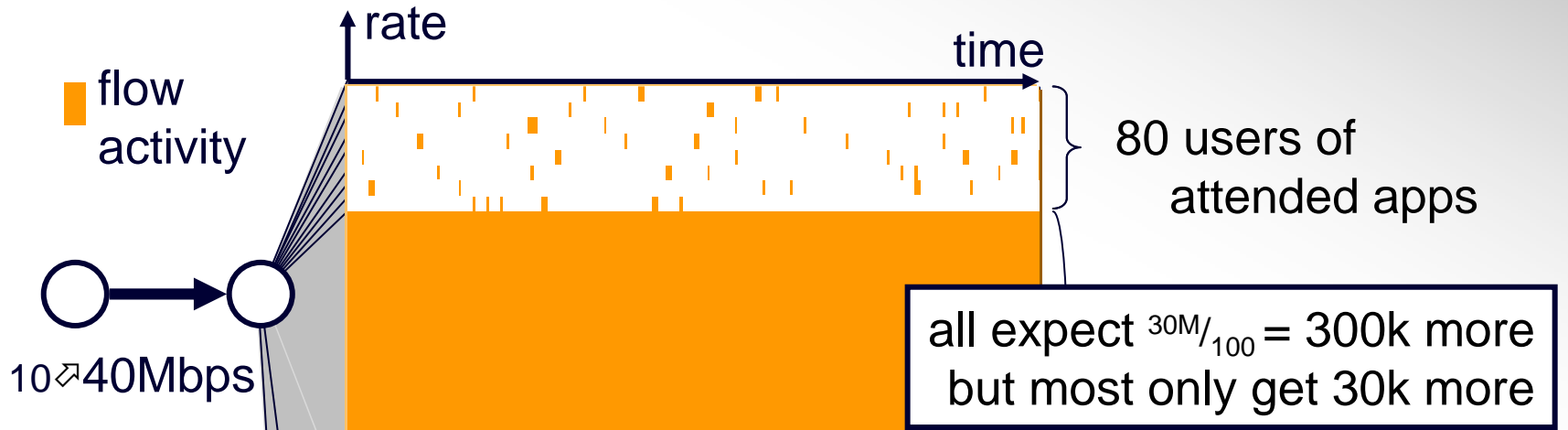
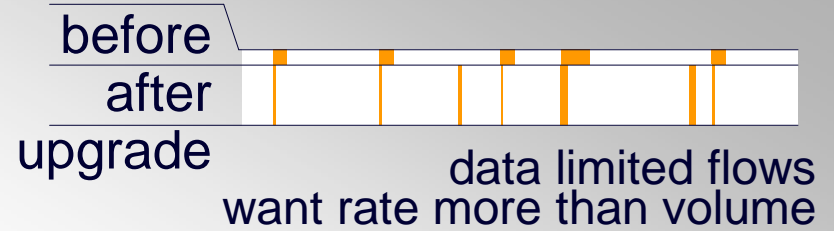
compounding activity factor & multiple flows

no-one is saying more volume is unfair
but volume accounting says it's fairer if heavier users get less rate during peak period



| usage type | no. of users | activity factor | ave.simul flows /user | TCP bit rate /user | vol/day (16hr) /user | traffic intensity /user |
|------------|--------------|-----------------|-----------------------|--------------------|----------------------|-------------------------|
| attended | 80 | 10% | 2 | 10kbps | 7.1MB | 1kbps |
| unattended | 20 | 100% | 100 | 500kbps | 3.6GB | 500kbps |
| | | | | x50 | x500 | x500 |

most users hardly benefit from bottleneck upgrade



2Mbps access each

| usage type | no. of users | activity factor | ave. simul flows /user | TCP bit rate /user | vol/day (16hr) /user | traffic intensity /user |
|------------|--------------|-----------------|------------------------|--------------------|----------------------|-------------------------|
| attended | 80 | 4% | 2 | 10 ↗ 40kbps | 12MB | 1 ↗ 1.6kbps |
| unattended | 20 | 100% | 100 | 0.5 ↗ 2Mbps | 14GB | 0.5 ↗ 2Mbps |

x50

x1250

20 users of unattended apps

volume accounting isn't the answer either

- fairer if heavy users get less bottleneck flow rate than light users
 - but heavy & light only defined by volume during 'the peak period'
 - effectively treats congestion very vaguely as
 - 0 everywhere off-peak
 - 1 everywhere on-peak
 - blind to whether the same volume causes extreme congestion or none

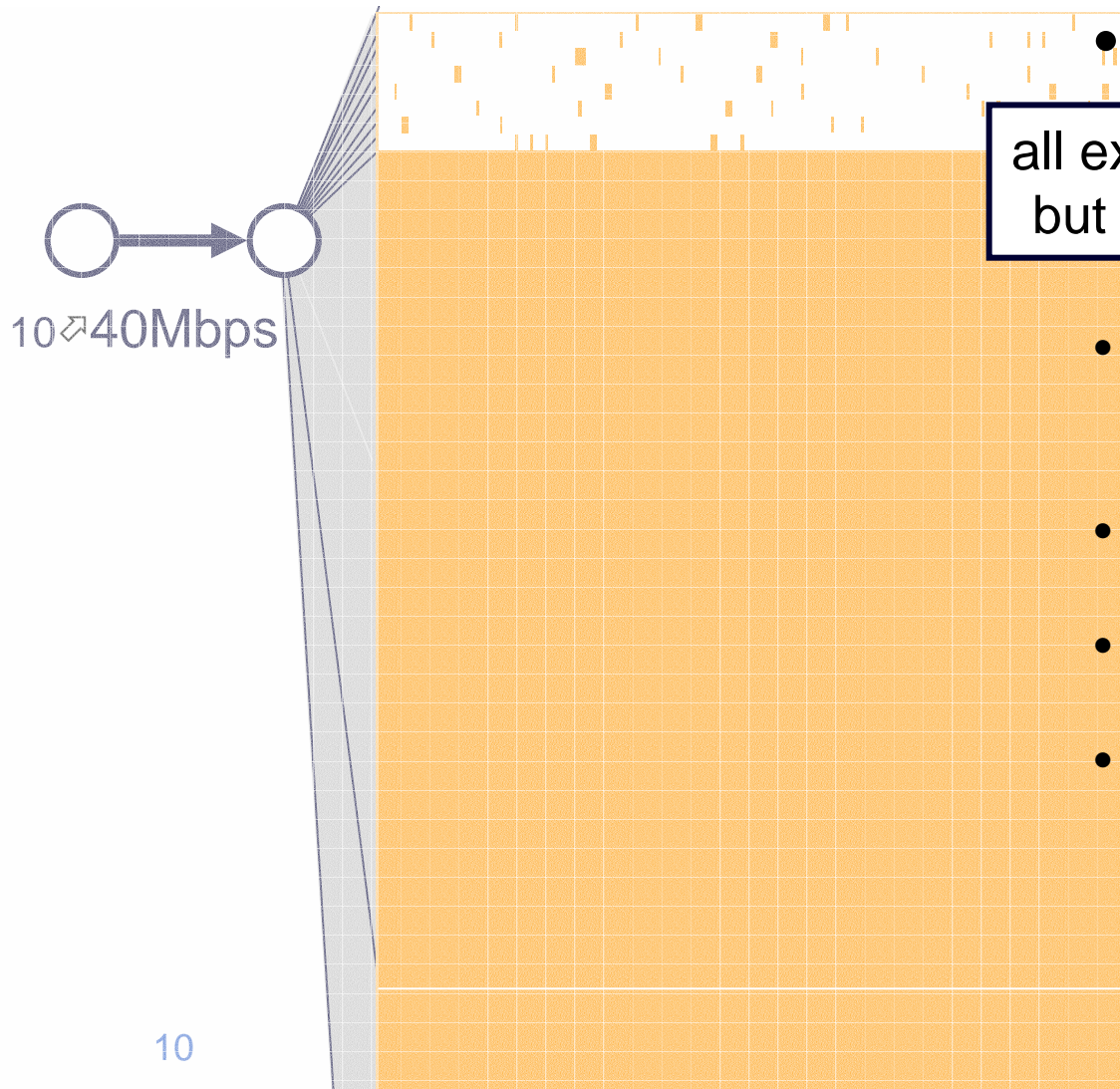
| degree of freedom | 'flow rate equality' | 'volume accounting' |
|----------------------|----------------------|---------------------|
| multiple flows | x | ✓ |
| activity factor | x | ✓ |
| congestion variation | ✓ | x |

- message so far: 2 worldviews both claim same goal (fairness)
 - each strong over part of the problem space
 - but incompatible: one wants equal, the other wants unequal flow rates

so what?

- fairness can't be such a problem, the Internet works
 - we all have enough most of the time, even if A has more than B
 - we like to think this is due to IETF protocols
 - next few slides cast doubt on this complacency

concrete consequence of unfairness #1 higher investment risk



- recall

all expect $30M/100 = 300k$ more
but most only get 30k more

- but ISP needs everyone to pay for 300k more
- if most users unhappy with ISP A's upgrade
- they will drift to ISP B who doesn't invest
- competitive ISPs will stop investing...

...but we still see enough investment

- main reasons
 - subsidies (e.g. Far East)
 - light users get 'enough' if more investment than they pay for
 - weak competition (e.g. US)
 - operators still investing because customers will cover the costs
 - throttling heavy users at peak times (e.g. Europe)
 - overriding TCP's rate allocation

concrete consequence of unfairness #2

trend towards bulk enforcement

- as access rates increase
 - attended apps leave access unused more of the time
 - anyone might as well fill the rest of their own access capacity
- operator choices:
 - a) either continue to provision sufficiently excessive shared capacity
 - b) or introduce tiered volume limits etc
- IETF needs to recognise & address the implications
 - bulk policing prevalent in best efforts architecture (cf. Diffserv)
 - e.g. should we distinguish a policer drop from a congestion drop?

concrete consequence of unfairness #3

networks making choices for users

- networks hit a problem once they start throttling
 - they could throttle all a heavy user's traffic indiscriminately
 - encourages the user to self-throttle least valued traffic
 - but many users have neither the software nor the expertise
- many networks *infer* what the user would do
 - using deep packet inspection (DPI) to identify apps
- even if intentions honourable
 - confusable with attempts to discriminate against certain apps
 - user's priorities are task-specific, not app-specific
 - customers understandably get upset when ISP guesses wrongly
- IETF needs to recognise & address the underlying need here
 - feature creep into network slows innovation (e2e principle)
 - better ways to fit traffic within limits (e.g. user/app-controlled endpoint s/w)

the problem

- IETF doesn't really decide fairness
 - whatever protocols *designed* to do, they are being *used* unfairly
- IETF can't really decide fairness
 - design-time body can't control run-time degrees of freedom
- IETF shouldn't decide fairness
 - shouldn't prejudge fair-use policy agreed between user & ISP
 - whether TCP, max-min, proportional or cost fairness

what does the IETF need to do?

- average rates – a run-time issue
 - introduce congestion accountability framework*
 - give principled effective fairness control to users, apps & operators
 - offer an evolvable alternative to current kludges (DPI)
 - coexist with null enforcement
- transport dynamics – the design-time issue
 - IETF/IRTF protocols can truly satisfy dynamic application requirements while minimising congestion
 - rather than not really meeting app reqs, by being over-constrained

* TBA (Lou Burness +)

working towards BoF, not just about fairness, but also congestion collapse & DDoS
re-ECN / re-feedback one proposed solution

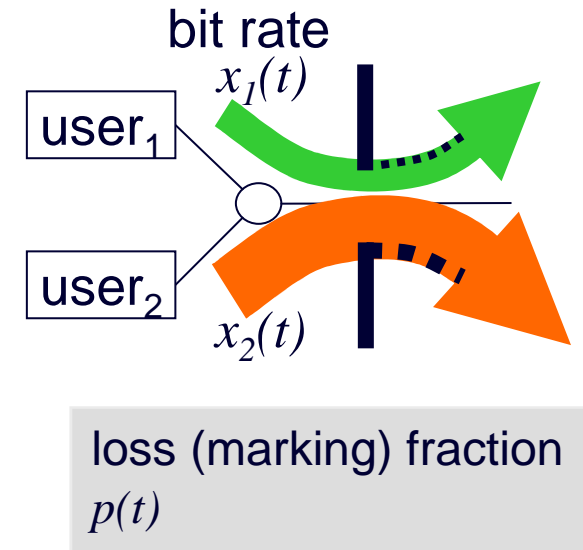
relaxing our transport design constraints

- currently we are trying to satisfy demanding app reqs
 - constrained by staying not 'much' more demanding than TCP
 - resulting protocols are 'over-constrained' and not app-developer's first choice
- once the big *average* rate fairness trade-offs move to run-time
- IETF/IRTF can judge which proposed transports better trade-off:
 - achieving the task effectively and
 - minimising unnecessary congestion to others during *dynamics*

- focus on the demanding *dynamics* questions:
 - when is a fast start fast enough? or too fast? [Limited slow start, etc]
 - how quickly should hi-speed transports allow in new flows? [HighSpeed TCP, FAST, etc]
 - how smooth can a transport be before it's effectively unresponsive? [TFRC, proprietary media players, etc]
 - what's the minimum congestion response of an aggregate? [PWE3, CAPWAP]

proposed core of solution congestion harm metric

- partial insight from volume accounting
- but rather than only counting bytes during peak
 - count bit rate *weighted* by congestion, over time
 - result is easy to measure per flow or per user
 - volume of bytes discarded (or ECN marked)
- termed **congestion volume**
- a precise instantaneous measure of harm, counted over time
 - a measure for fairness over any timescale
 - and a precise measure of harm during dynamics



summary

shift IETF focus from fairness to accountability

| | design-time | run-time |
|--|--|---|
| problem | IETF doesn't, can't and shouldn't decide fairness | users, apps & operators actually control fairness |
| solution process | IETF's role: enable accountability for congestion | users, apps & operators can (optionally) make principled fairness choices |
| | IETF/IRTF can truly meet dynamic app req's and minimise congestion | IETF protocols become first choice for demanding apps 😊 |
| best metric: congestion volume | | |

- problems will only get worse – driven by access rate increases

we don't have to decide
fairness ourselves

[<draft-briscoe-tsvwg-relax-fairness-00.txt>](#)

Q&A



context

3. a protocol solution: re-ECN <[draft-briscoe-tsvwg-re-ecn-04.txt](#)>
 - on hold while build consensus on the problem & requirements
 - other solutions welcome
0. dismantling flow rate fairness <[draft-briscoe-tsvarea-fairness-02.pdf](#)>
 - too polemical for IETF consensus
 - let this draft die – archived on my Web site and ACM CCR paper
1. the problem <[draft-briscoe-tsvwg-relax-fairness-00.txt](#)>
 - IETF doesn't decide fairness – this talk
2. solution requirements <draft-burness-tsvwg-...>
 - TBA

not pushing technical solution(s) at steps 1 & 2

- aimed more towards a 'congestion accountability' BoF

typical p2p file-sharing apps

- 105-114 active TCP connections altogether

1 of 3 torrents shown

- ~45 TCPs per torrent
- but ~40/torrent active

environment

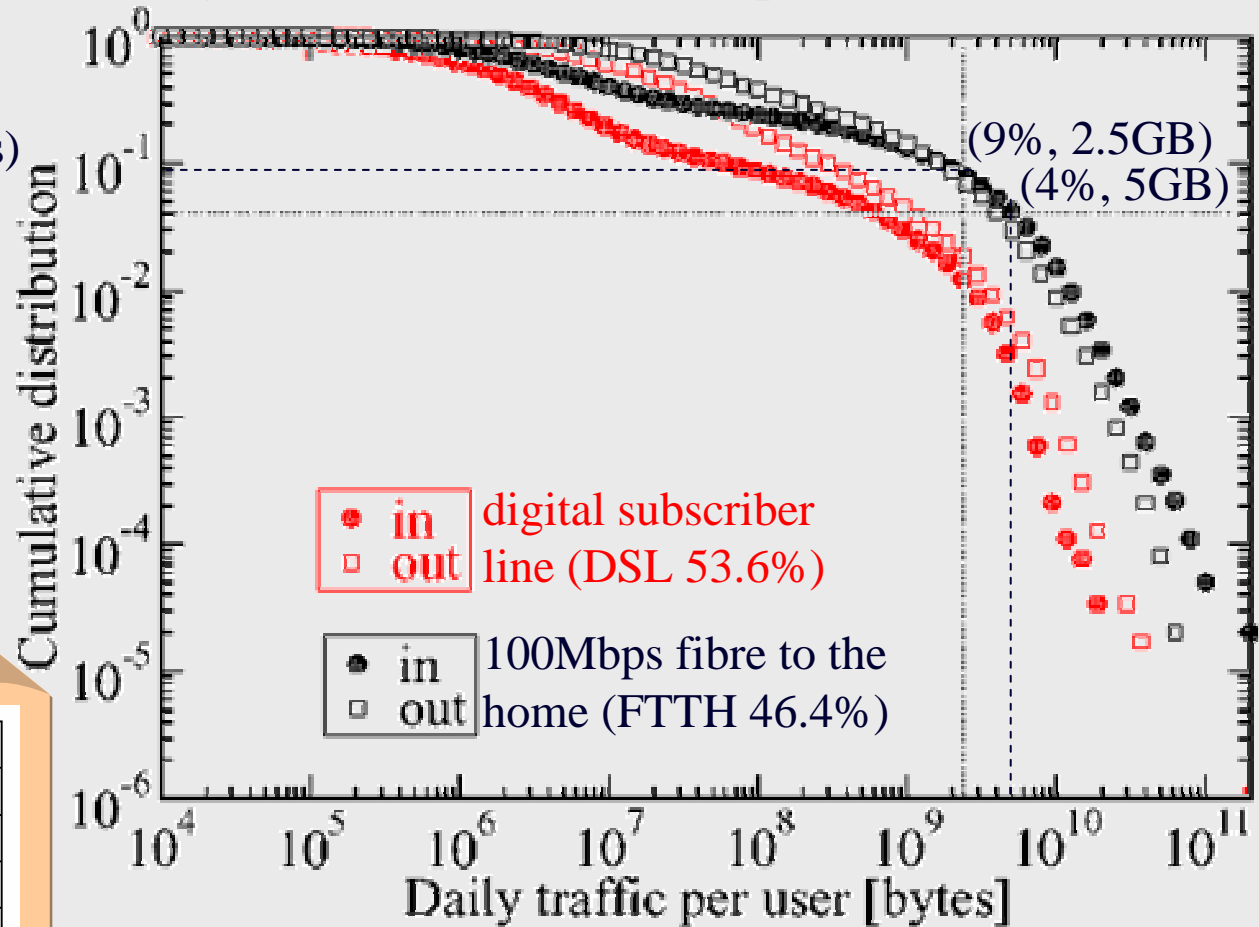
Azureus BitTorrent app
ADSL+ 448kb upstream
OS: Windows XP Pro SP2

| IP | Client | T | Pieces | % | D... | Up Speed | State | Encryption | Down | Up | I |
|--------------|-----------------|-----|--------|--------|-----------|----------|-------------------|------------|----------|---------|-----|
| 78.86.8.10 | Azureus 3.0.2.2 | L | | 100.0% | 14.5 kB/s | 44 B/s | Fully established | RC4-160 | 6.87 MB | 25.8 kB | |
| 76.65.28.192 | µTorrent 1.7.5 | R | | 100.0% | 11.1 kB/s | 20 B/s | Fully established | None | 10.52 MB | 14.6 kB | |
| ... | ... | ... | ... | ... | ... | ... | ... | ... | ... | ... | ... |

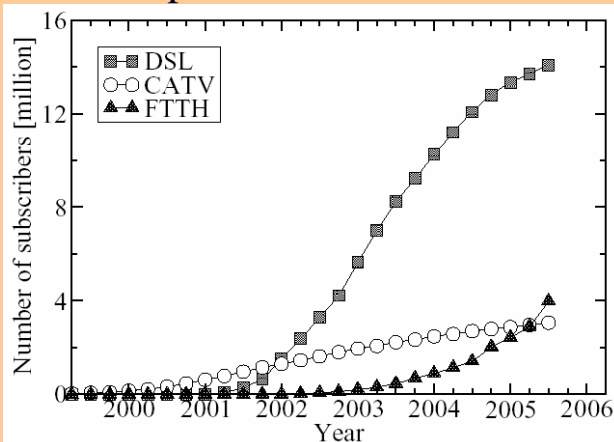
access growth just gets filled

Distribution of customers' daily traffic into & out of a Japanese ISP (Feb 2005)

(5GB/day equivalent to
0.46Mbps if continuous)



Changing technology shares of Japanese access market



Courtesy of Kenjiro Cho et al
The Impact and Implications of the Growth
in Residential User-to-User Traffic, SIGCOMM'06

concrete consequence of unfairness #4 starvation during anomalies & emergencies

- fairness concerns become acute during stress
 - more traffic or less capacity than expected
- if fairness decided at run-time
 - common policy probably 'you get what you paid for'
- concern: unsavoury for emergencies
 - all flows should make some progress, not just the rich
- agree with concern, but current approach not right
 - video downloads get 50x rate of emergency messages?*
- policy decisions for users, ISPs, regulators, not IETF
 - e.g. ISP might freeze paying to override congestion limits

* Henchung earthquake, 26 Dec '06, see I-D

accountability metric

congestion volume

precisely measures instantaneous harm from flow rate dynamics rather than just average flow rate

