



QoS interconnect best without effort

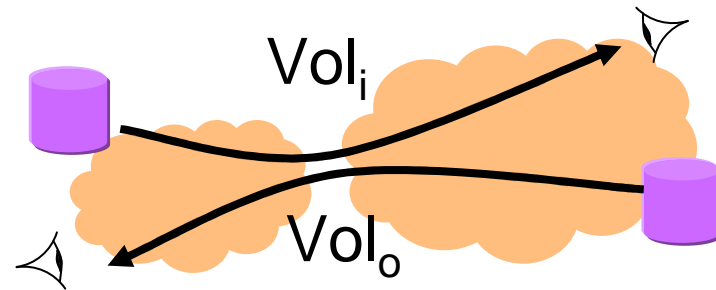
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Chief Researcher
BT
Sep 2009

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supported by the European Community
www.trilogy-project.org

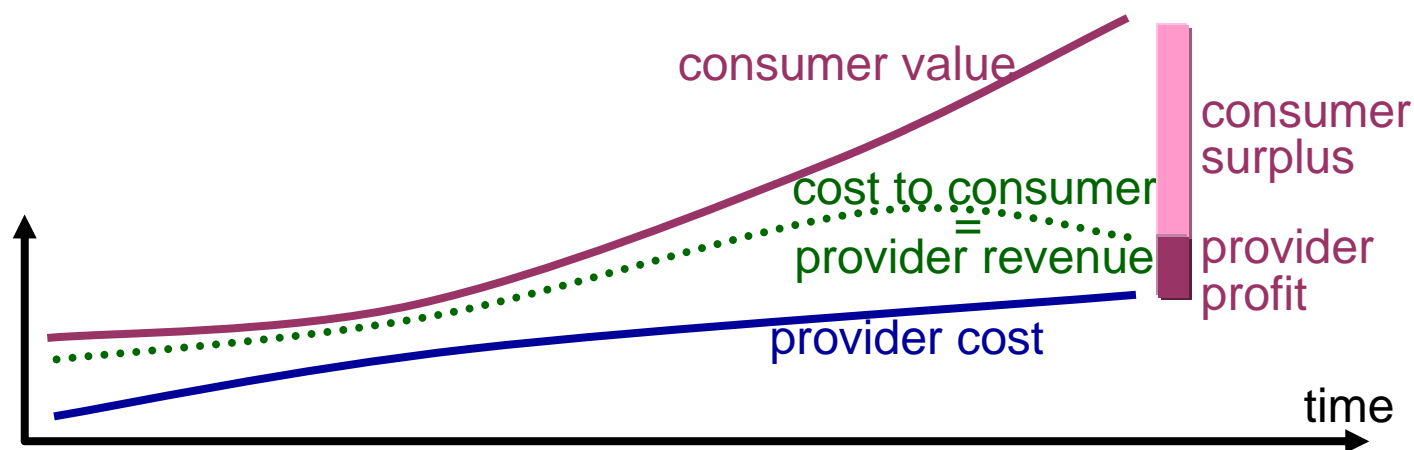




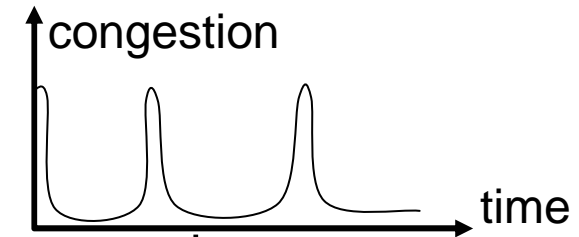
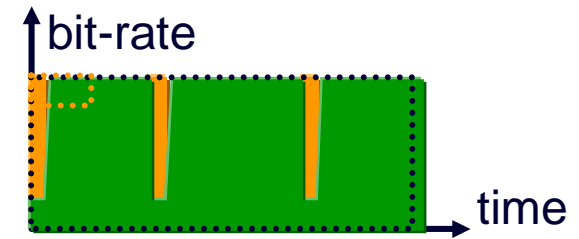
both value *and* cost



- industry contractual metrics are largely value-based
 - e.g. advertised routes, volume ratio
 - even a CEO should understand both value *and* cost
- competitive market drives revenues down towards provider's marginal cost
 - those who understand marginal costs will succeed

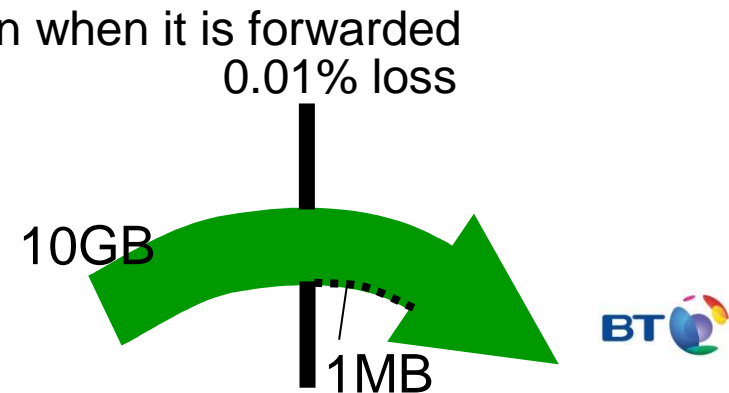
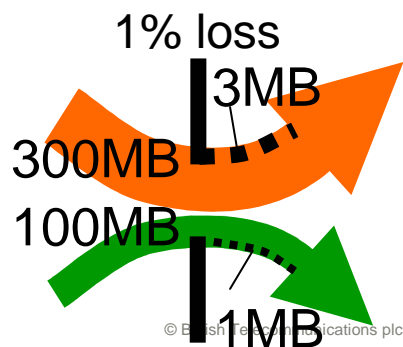


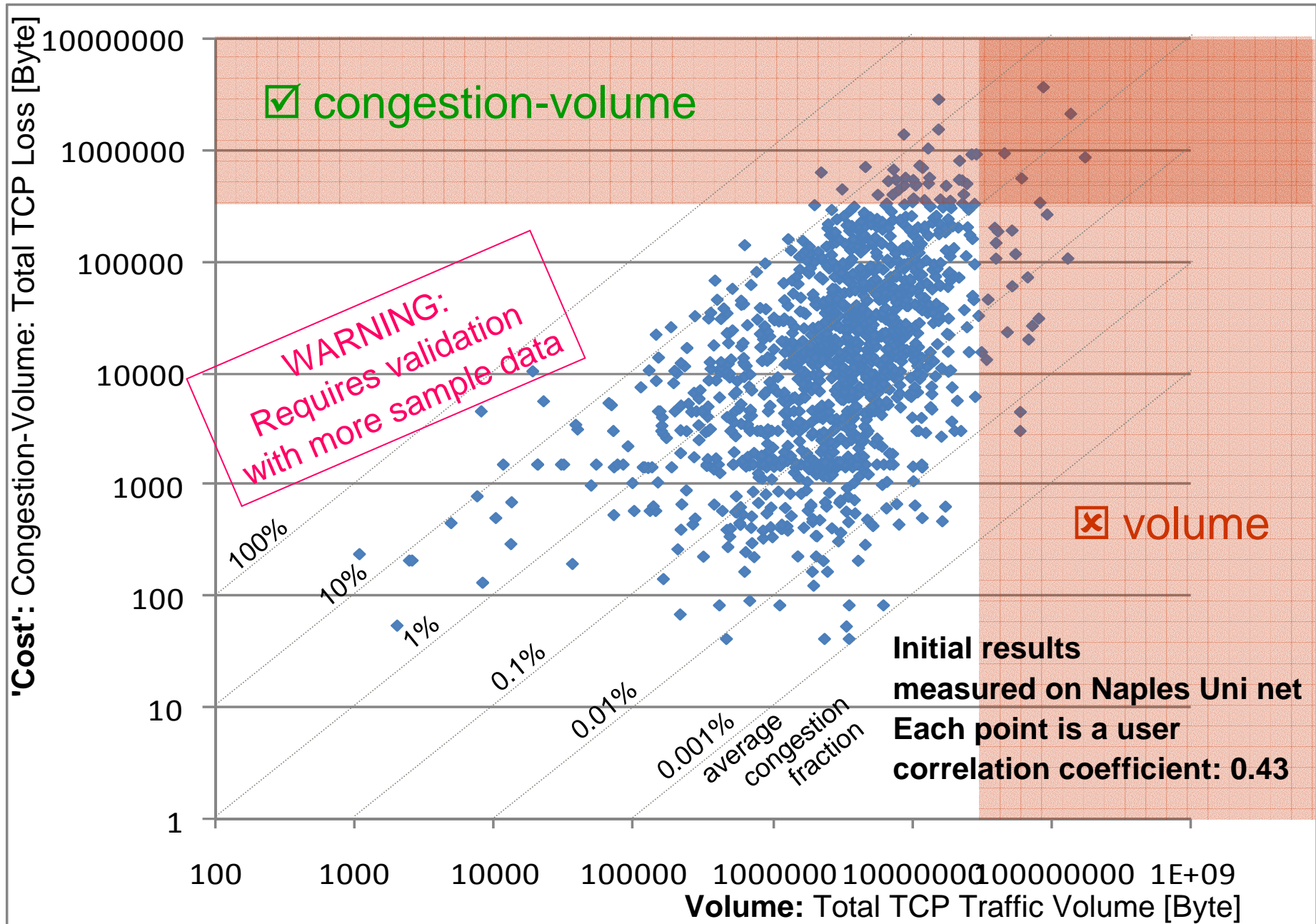
marginal cost of network usage?



- volume is NOT a good measure
- green user yields whenever detects high congestion
 - very high volume but very low cost to others
 - e.g. LEDBAT (BitTorrent's low extra delay background transport) or weighted TCP
- by counting volume, ISPs kill nice behaviour
 - not just file transfers, e.g. congestion-sensitive video codec transfers >100% more videos thru same capacity (same MoS)
- correct measure: congestion-volume

- volume weighted by congestion when it is forwarded
- easily measured by a host
- bytes sent x loss fraction = bytes lost

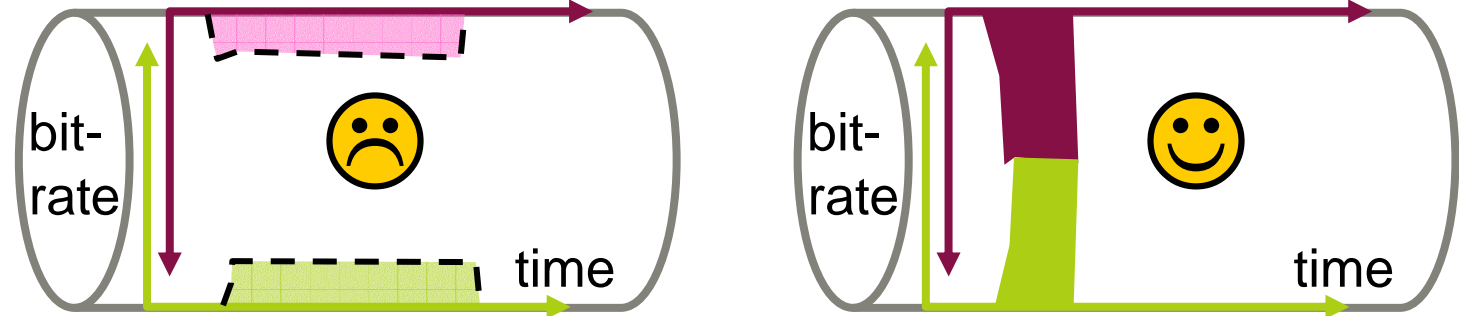




congestion is not evil

congestion signals are healthy

- no congestion across whole path \Rightarrow feeble transport protocol
 - to complete ASAP, transfers should sense path bottleneck & fill it



the trick

congestion signal *without* impairment

- explicit congestion notification (ECN)
 - update to IP in 2001: mark more packets as queue builds
- then tiny queuing delay and tiny tiny loss for all traffic
- no need to avoid congestion (whether core, access or borders) to prevent impairment

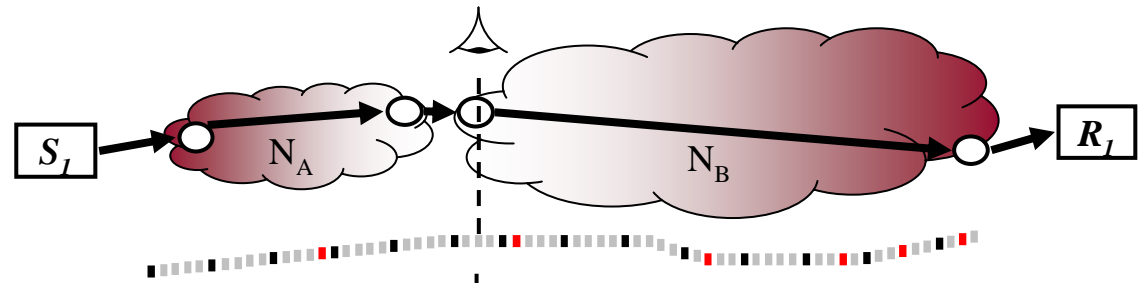
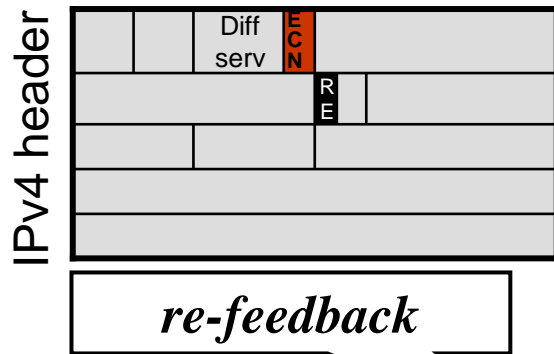
congestion exposure

- by Internet design, endpoints detect & handle losses
 - v hard for networks to see losses (marginal costs)
- proposed IETF working group: “congestion exposure”
 - protocol for sender to mark IP headers to expose congestion
 - to measure traffic cost as easily as we measure volume
 - just count volume of marked packets in aggregate
 - >40 offers of help just in the last fortnight
- named re-ECN (re-inserted ECN)
 - builds on explicit congestion notification (ECN [RFC3168])

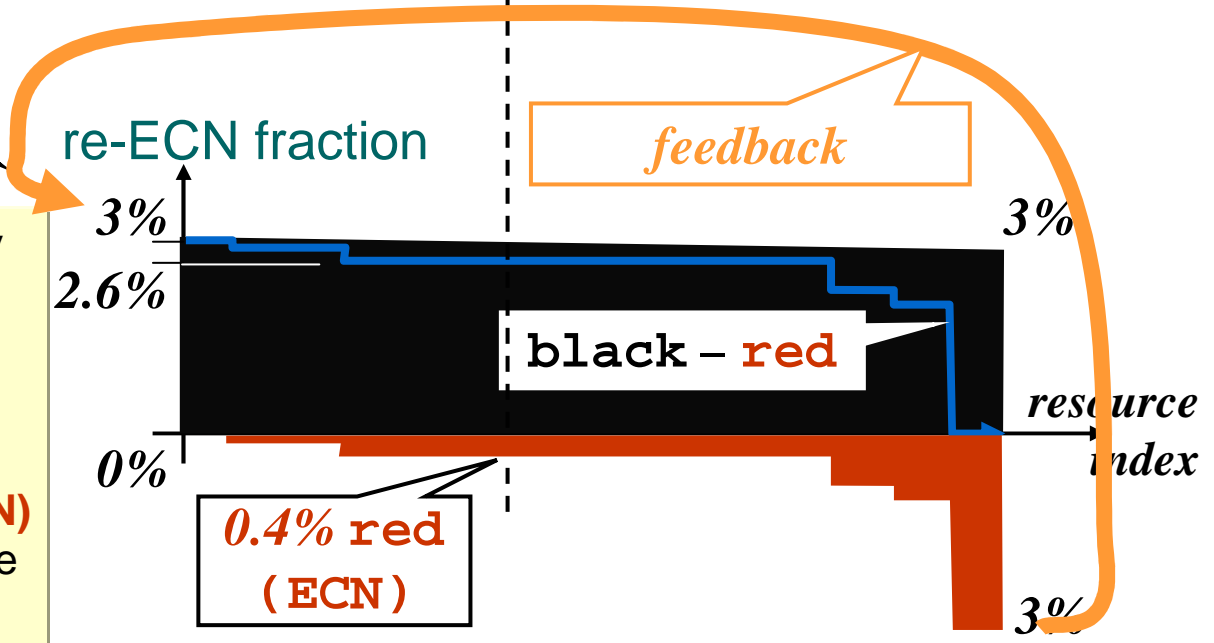
I E T F[®]

congestion exposure with ECN & re-ECN

measurable upstream, downstream and path congestion



- sender re-inserts feedback by marking packets **black**
- at any point on path, diff betw fractions of **black** & **red** bytes is downstream congestion
- **forwarding unchanged (ECN)**
- **black** marking e2e but visible at net layer for accountability

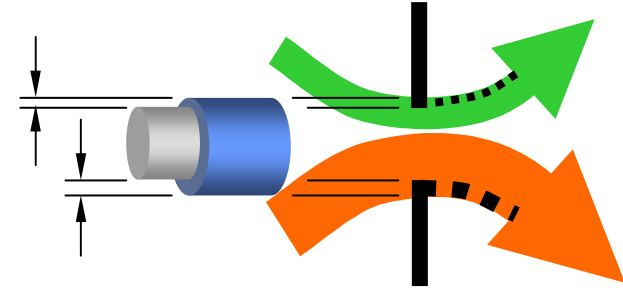


congestion-volume metric dual demand & supply role

- a resource accountability metric
 1. of customers to ISPs (too much traffic)
 2. and ISPs to customers (too little capacity)

- 1. cost to other users of my traffic
- 2. the marginal cost of upgrading equipment
 - so it wouldn't have been congested

- competitive market matches 1 & 2



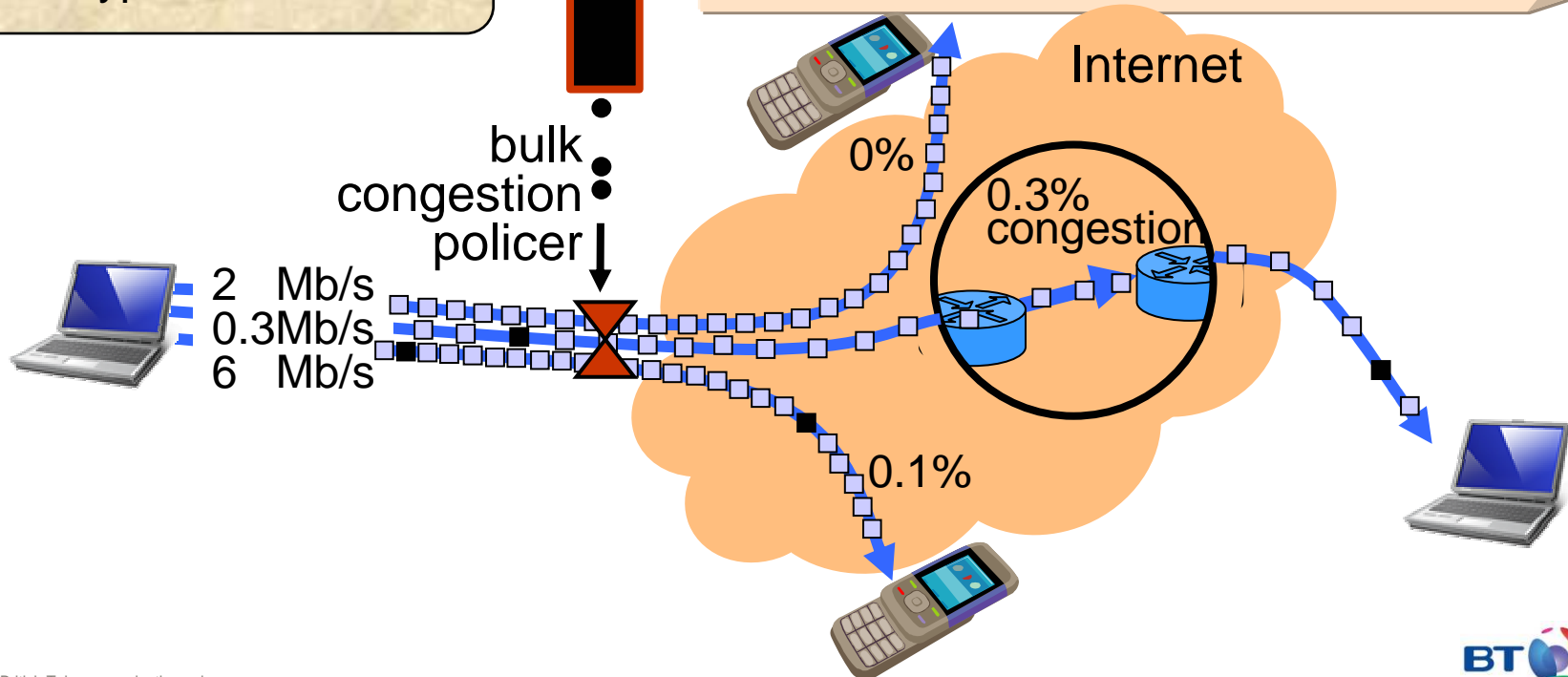
example consumer use of exposed congestion fee can stay flat

Acceptable Use Policy

'congestion-volume'
allowance: 1GB/month

Allows ~70GB per day of
data in typical conditions

- only throttles congestion-causing traffic when your contribution to congestion EVERYwhere in the Internet exceeds your allowance
- side-effect: mitigates and reveals distributed denial of service



reveals congestion dumped
into rest of Internet

legend:

re-ECN
downstream
congestion
marking [%]

bit rate

area =
instantaneous
downstream
congestion
volume

N_B could be an IX

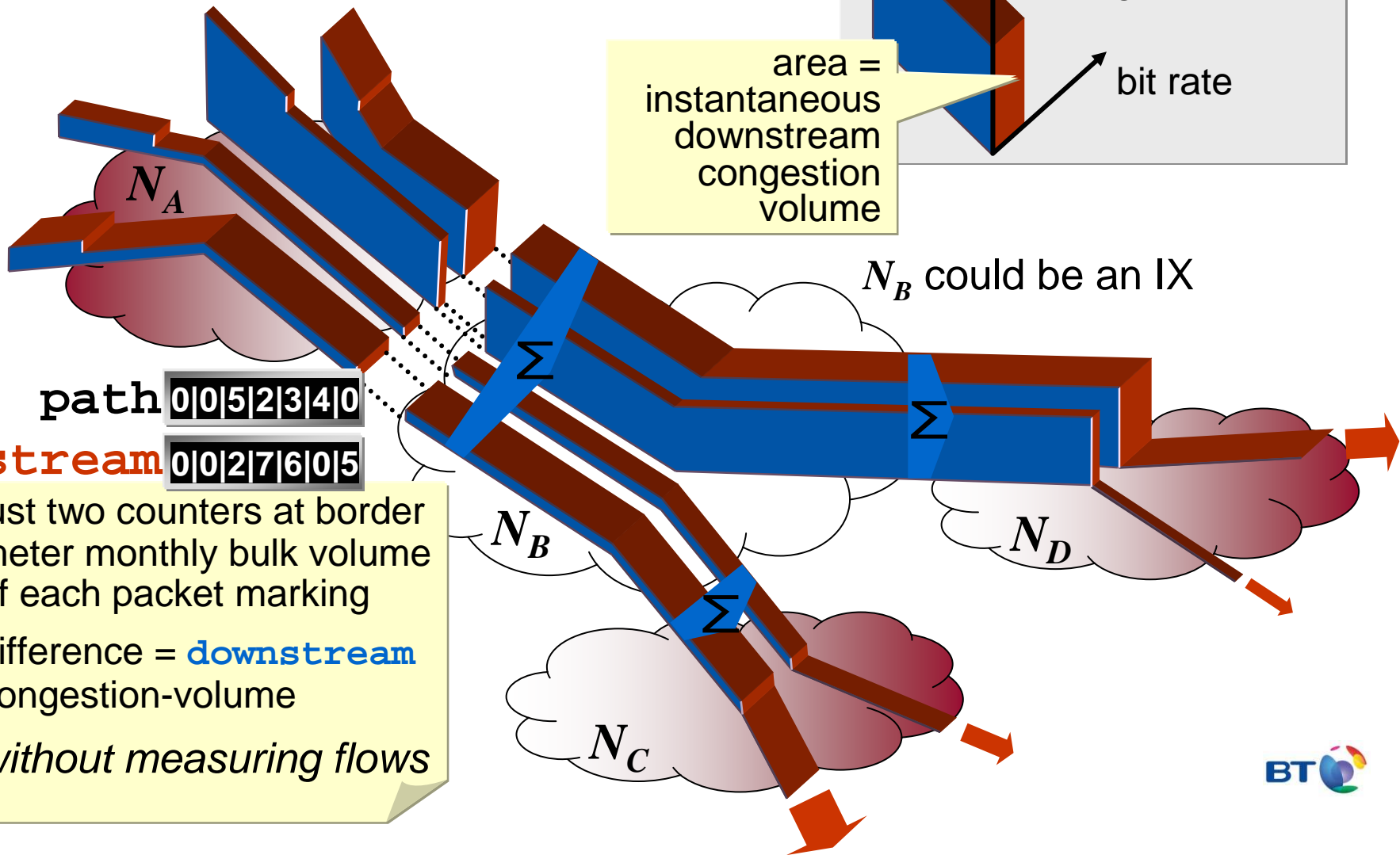
path 0|0|5|2|3|4|0

upstream 0|0|2|7|6|0|5

just two counters at border
meter monthly bulk volume
of each packet marking

difference = downstream
congestion-volume

without measuring flows



I'm a conservative, get me out of here!

- if we don't listen to the economics, we're all dead
 - shift from value-based to cost-based is unstoppable
 - competition
 - bit transport needs to be viable on its own

(another talk)

- as cost pressures grow
- existing capacity sharing methods feed an arms race
 - TCP doesn't share capacity fairly by any means
 - recent unanimous consensus in IETF Transport Area
 - ISPs have quietly been fighting TCP with piecemeal tools
 - WFQ, volume capping, deep packet inspection
- with congestion in IP header, wouldn't need to look deeper

best without effort

- did you notice the interconnected QoS mechanism?
 - *endpoints* ensure tiny queuing delay & loss for all traffic
 - if your app wants more bit-rate, it just goes faster
 - effects seen in bulk metric at every border (for SLAs, AUPs)
- simple – and all the right support for operations
- the invisible hand of the market
 - favours ISPs that get their customers to manage their traffic in everyone else's best interests
- incentives to cooperate across Internet value chain
 - content industry, CDNs, app & OS authors, network wholesalers & retailers, Internet companies, end-customers, business, residential
- if you want this, vote early and vote often!
 - re-ecn@ietf.org list
 - IETF, Hiroshima, Nov'09



more info...

- White paper – the whole story in 7pp
 - [Internet: Fairer is Faster](#), Bob Briscoe (BT), BT White Paper TR-CXR9-2009-001 (May 2009)
 - an abridged version of this article appeared in IEEE Spectrum, Dec 2008
- Inevitability of policing
 - The Broadband Incentives Problem, Broadband Working Group, MIT, BT, Cisco, Comcast, Deutsche Telekom / T-Mobile, France Telecom, Intel, Motorola, Nokia, Nortel (May '05 & follow-up Jul '06) <cfp.mit.edu>
- Stats on p2p usage across 7 Japanese ISPs with high FTTH penetration
 - Kenjiro Cho et al, "The Impact and Implications of the Growth in Residential User-to-User Traffic", In Proc ACM SIGCOMM (Oct '06)
- Slaying myths about fair sharing of capacity
 - Bob Briscoe, "[Flow Rate Fairness: Dismantling a Religion](#)" ACM Computer Communications Review 37(2) 63-74 (Apr 2007)
- How wrong Internet capacity sharing is and why it's causing an arms race
 - Bob Briscoe et al, "[Problem Statement: Transport Protocols Don't Have To Do Fairness](#)", IETF Internet Draft (Jul 2008)
- Understanding why QoS interconnect is better understood as a congestion issue
 - Bob Briscoe and Steve Rudkin "[Commercial Models for IP Quality of Service Interconnect](#)" BT Technology Journal 23 (2) pp. 171--195 (April, 2005)
- Re-architecting the Internet:
 - The [Trilogy](#) project
- Re-ECN & re-feedback project page:
 - <<http://bobbriscoe.net/projects/refb/>>
 - <trac.tools.ietf.org/area/tsv/trac/wiki/re-ECN>

best without effort
QoS interconnection

Q&A...



problems using congestion in contracts

| | 1. loss | 2. ECN | 3. re-ECN |
|---|---------|--------|-----------|
| can't justify selling an impairment | ☹ | ☺ | ☺ |
| absence of packets is not a contractible metric | ☹ | ☺ | ☺ |
| congestion is outside a customer's control | ☹ | ☹ | ☺ |
| customers don't like variable charges | ☹ | ☹ | ☺ |
| congestion is not an intuitive contractual metric | ☹ | ☹ | ☹ |

1. **loss**: used to signal congestion since the Internet's inception

- computers detect congestion by detecting gaps in the sequence of packets
- computers can hide these gaps from the network with encryption

2. **explicit congestion notification (ECN)**: standardised into TCP/IP in 2001

- approaching congestion, a link marks an increasing fraction of packets
- implemented in Windows Vista (but off by default) and Linux, and IP routers (off by default)

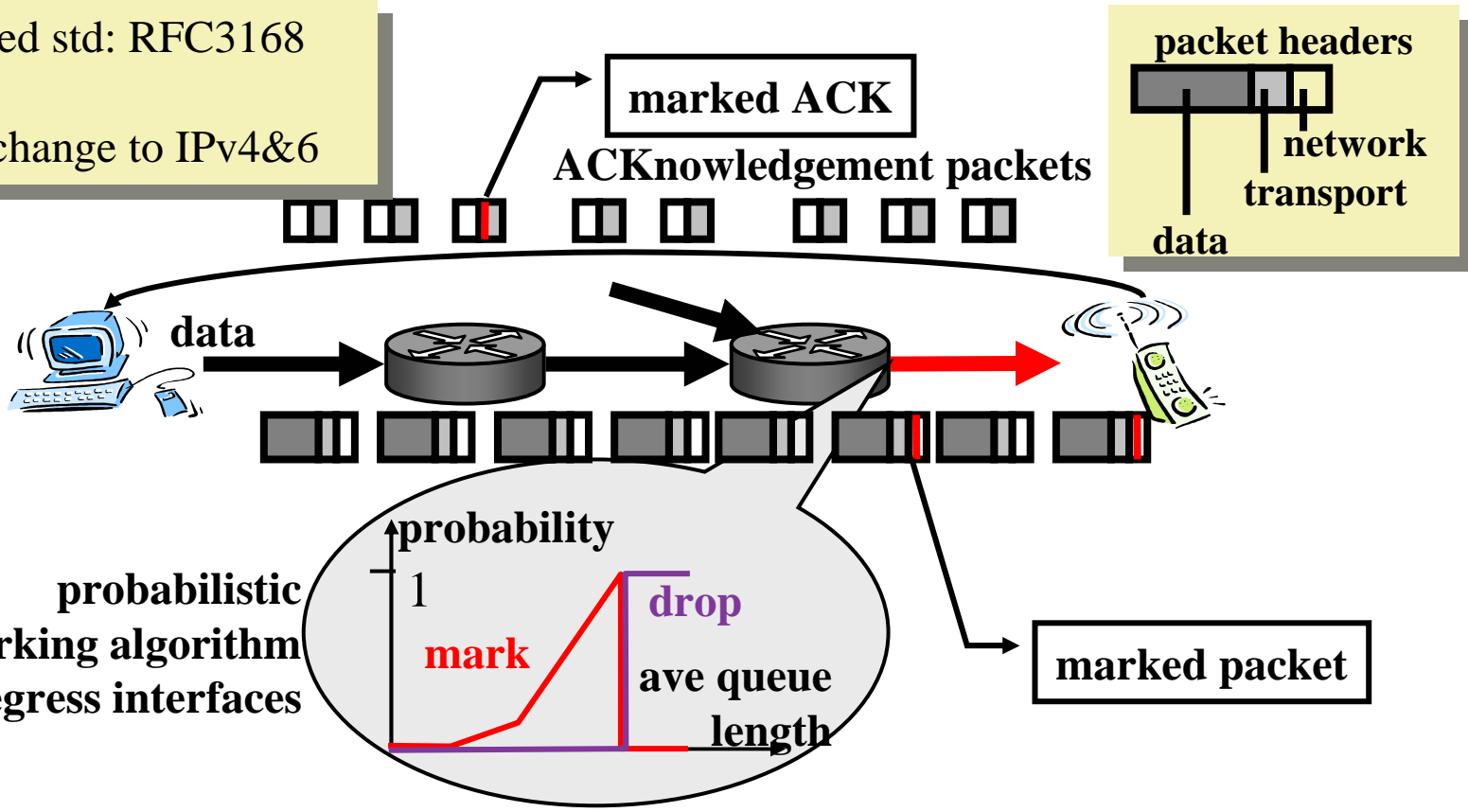


3. **re-inserted ECN (re-ECN)**: standards proposal since 2005

- packet delivery conditional on sender declaring expected congestion
- uses ECN equipment in the network unchanged

explicit congestion notification (ECN)

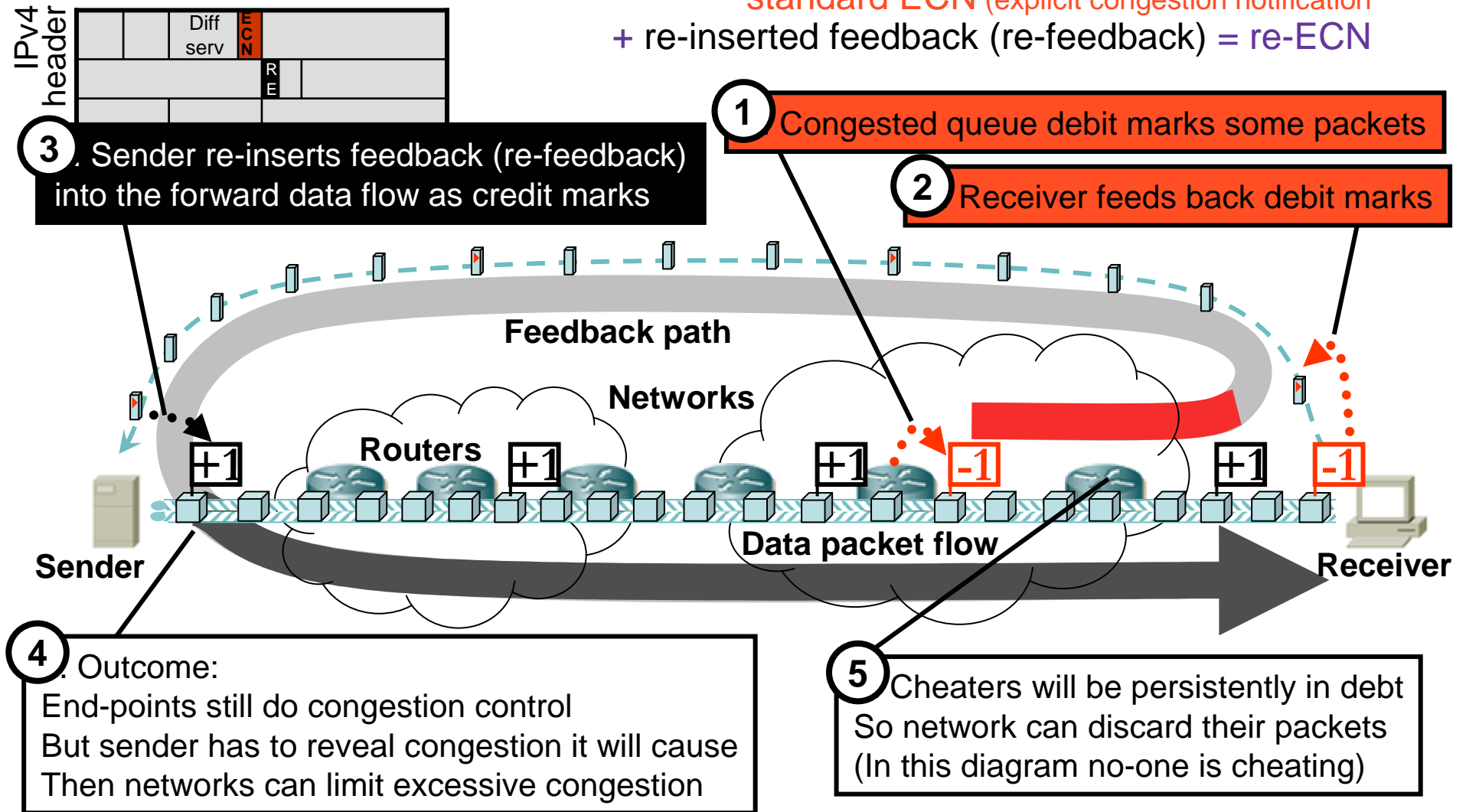
IETF proposed std: RFC3168
 Sep 2001
 most recent change to IPv4&6



| | | | |
|-----------|--|--------------------------|-------|
| 00: | Not ECN Capable Transport (ECT) | 0 | 5 6 7 |
| 01 or 10: | ECN Capable Transport - no Congestion Experienced (sender initialises) | DSCP | |
| 11: | ECN Capable Transport - and Congestion Experienced (CE) | ECN | |
| | | bits 6 & 7 of IP DS byte | |

congestion exposure in one bit

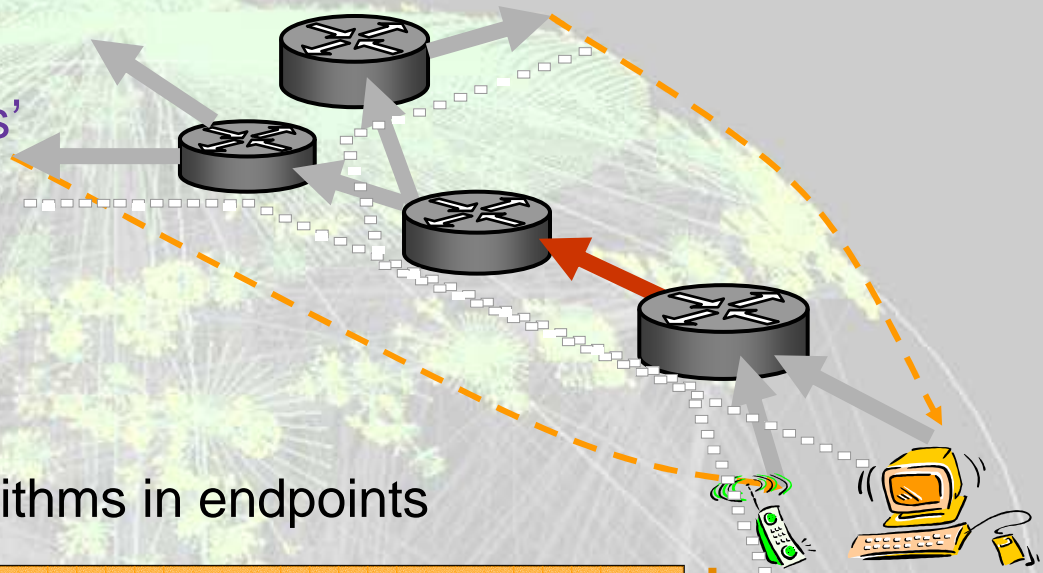
standard ECN (explicit congestion notification)
+ re-inserted feedback (re-feedback) = re-ECN



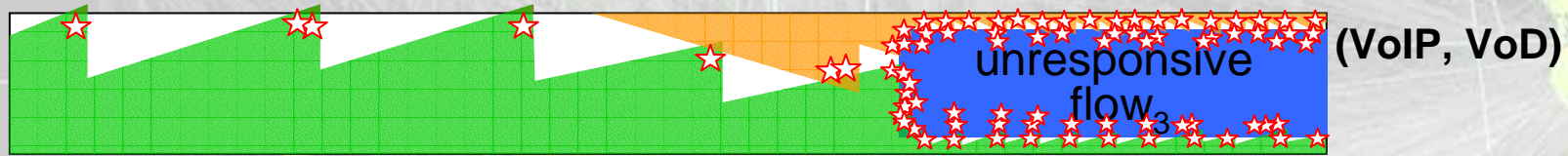
main steps to deploy re-feedback / re-ECN

- network
 - turn on explicit congestion notification in routers (already available)
 - deploy simple active policing functions at customer interfaces around participating networks
 - passive metering functions at inter-domain borders
- terminal devices
 - (minor) addition to TCP/IP stack of sending device
 - or sender proxy in network
- customer contracts
 - include congestion cap
- oh, and first we have to update the IP standard
 - started process in Autumn 2005
 - using last available bit in the IPv4 packet header
 - proposal for new working group, Nov 2009 IETF

how Internet sharing 'works' TCP-friendliness

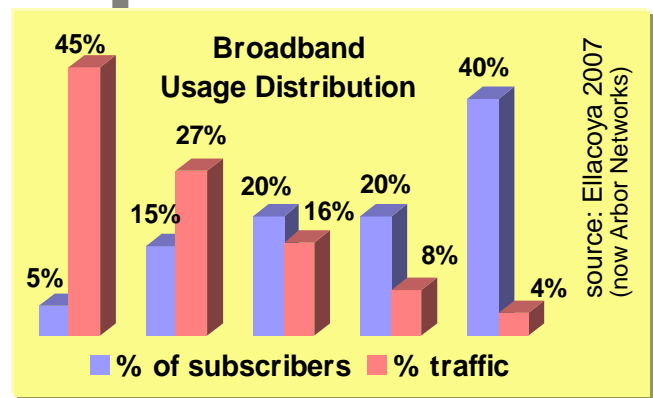
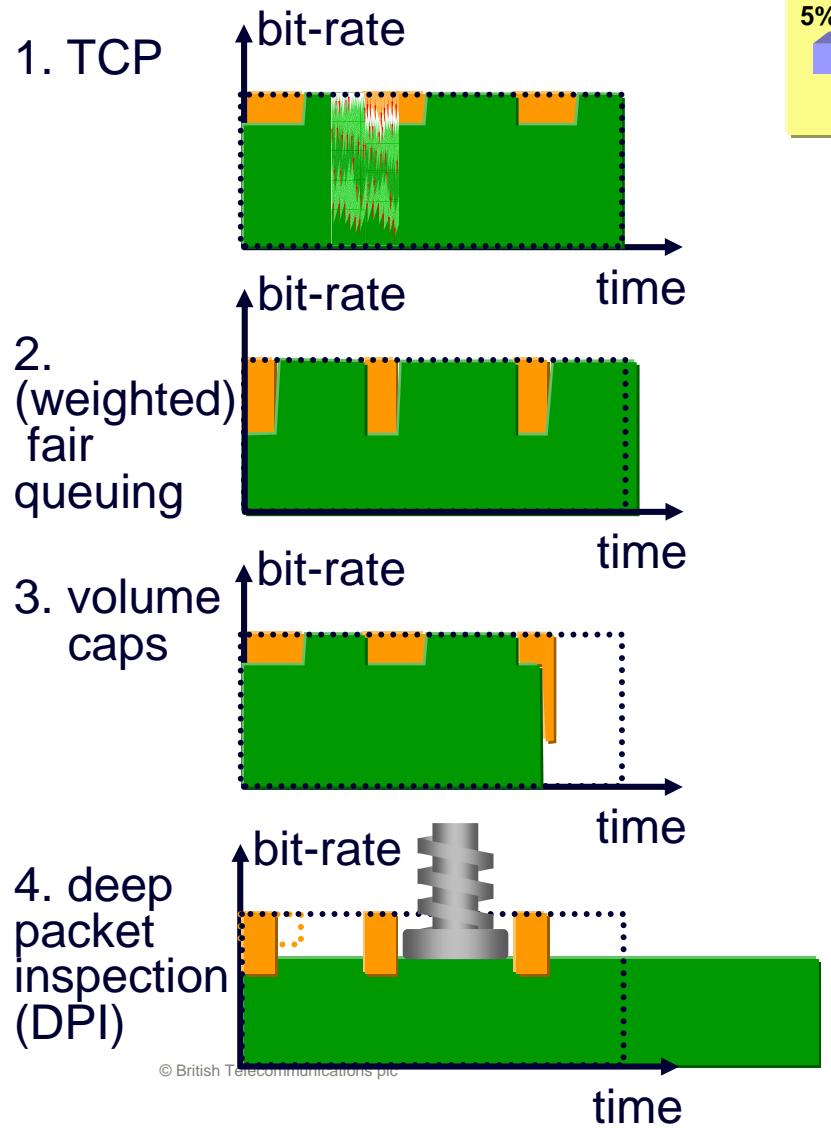


- endemic congestion
- voluntarily restraint by algorithms in endpoints



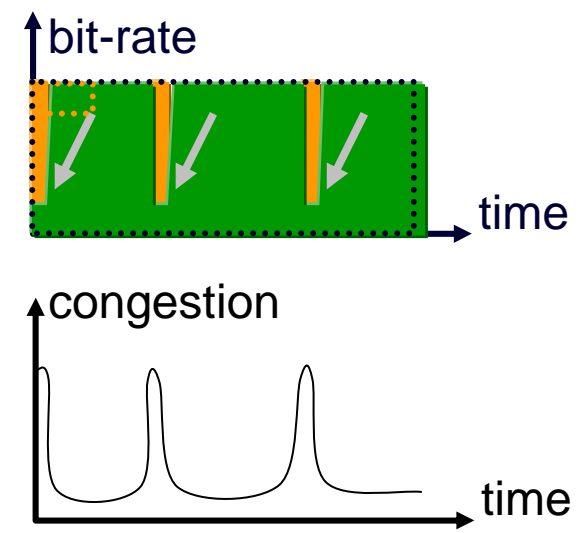
- or for much longer than anyone else (p2p file-sharing x200)
- net effect of both (p2p: x1,000-20,000 higher traffic intensity)

none of these
harness end-system flexibility



simpler
& better...

weighted
TCP
sharing



- light usage can go much faster
- hardly affects completion time of heavy usage

NOTE: weighted sharing doesn't imply differentiated network service

- just weighted aggressiveness of end-system's rate response to congestion



congestion competition – inter-domain routing

- if congestion \rightarrow profit for a network, why not fake it?
 - upstream networks will route round more highly congested paths
 - N_A can see relative costs of paths to R_1 thru N_B & N_C
- the issue of monopoly paths
 - incentivise new provision
 - as long as competitive physical layer (access regulation), no problem in network layer

