

# Byte and Packet Congestion Notification

[draft-briscoe-tsvwg-byte-pkt-mark-01.txt](#)

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IETF-70 tsvwg Dec 2007



# updated individual draft

- Byte and Packet Congestion Notification
  - **updated draft:** [draft-briscoe-tsvwg-byte-pkt-mark-01.txt](#)
  - **intended status:** informational
  - **immediate intent:** move to WG item

## reminder (exec summary)

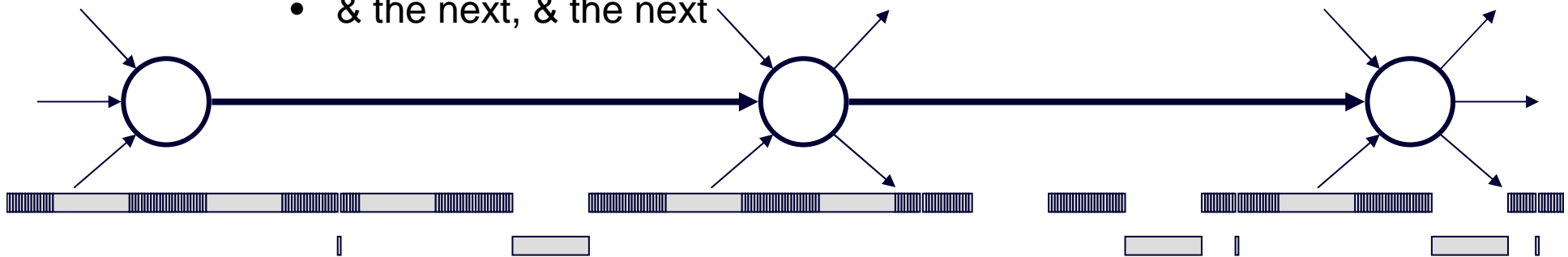
- in any AQM  
propose **SHOULD NOT** give smaller packets preferential treatment
- adjust for byte-size when transport reads **NOT** when network writes

Terminology: RED's 'byte mode queue measurement' (often called just 'byte mode') is OK, only 'byte mode packet drop' deprecated

**NOTE:** don't turn off RED completely: drop-tail is as bad or worse

## favouring small packets, main change: DoS vulnerability

- small packet attacks push out larger packets
  - leaving most small packets to attack the next queue
  - & the next, & the next



- DoS vulnerability similar to that of drop tail queues
- AQM was partly about not locking-out large packets\*
  - shouldn't add lock-out back again in the AQM algorithm

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\* not stated and not a motivation according to at least one author (Floyd)

# other changes

- emphasised equal applicability to any AQM and to drop or ECN
  - e.g. PCN, RED (with drop or ECN)
- restructured
  - pulled main recommendations together into the conclusions
  - moved a couple of lumps of text to appendices
- fixed for (Floyd's) original motivations for RED's byte-mode drop
  - protecting SYN's & pure ACKs more than equalising small segment TCPs
- added more examples of preferable transport approaches
  - tcpm-ecnsyn & tcpm-ackcc added to TFRC-SP etc
- updated survey data (but no change since IETF-69 slides)
- clarification & update throughout
- full diff at <[www.cs.ucl.ac.uk/staff/B.Briscoe/pubs.html#byte-pkt-mark](http://www.cs.ucl.ac.uk/staff/B.Briscoe/pubs.html#byte-pkt-mark)>

# thoughts for next draft

- long off-list discussions haven't resolved differences, but could
- main points in favour of size-dependent drop:
  - control packets tend to be small (e.g. SYNs, pure ACKs)
    - so less drop of small packets gives performance win
  - already have mix of size-dependent (drop-tail) and size-independent drop
    - so doesn't reduce complexity by only having size-independent
  - apps have other (OS) incentives not to use small packets
- main points in favour of size-independent drop
  - not all small pkts are control, so favouring all smallness creates unintended consequences
  - the more size-independent AQM, the less transport uncertainty over queue behaviours
  - mustn't provide incentives for new transports to use small data pkts
- possible ways forward
  - focus only on PCN?
  - but still mileage in reaching consensus on RED too



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## Q&A



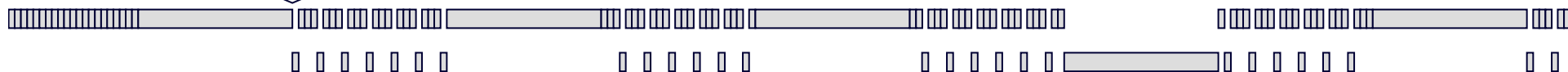
# example: comparing each RED mode

## simple packet streams (no congestion response)

RED  
packet-mode  
packet drop

- same drop probability for any packet
- universally deployed
- propose: **SHOULD**

	1500B pkts	60B pkts
input	1Mbps	1Mbps
drop prob.	25%	25%
output	750kbps	750kbps



RED  
byte-mode  
packet drop

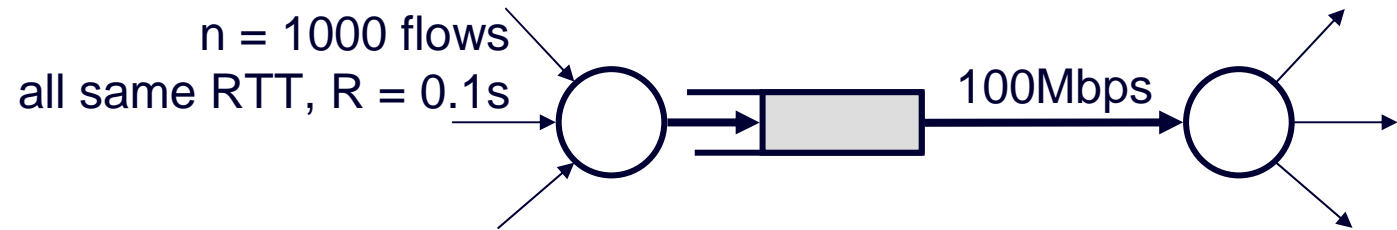
- lower drop probability for smaller packets
- 'RED' RFC2309 (sort of) recommends
- propose: **SHOULD NOT**

	1500B pkts	60B pkts
input	1Mbps	1Mbps
drop prob.	25%	1%
output	750kbps	990kbps





# proof



- proof strategy
  - fix the amount of congestion in flight, then consider how much notification needed
- Imagine aggregate overload of  $103$  Mbps in flight (for  $1$  RTT)
- arrives at queue running at its desired operating point  
 $\Rightarrow 3$  Mbps for  $0.1$  s =  $300$  kb =  $37.5$  kB to discard (or mark)
- If all the traffic is in packets of size  $s$  [B],  $e$  packets need to be lost

s	e
1500B	25
60B	625
9000B	4

- to lose this excess
  - if queue reduces drop rate for smaller packets
    - if all packets are small, queue will have to be longer
- to notify the same congestion queue shouldn't have to change its length
- so transports need to respond more strongly to larger missing packets
- $\therefore$  TCP's insensitivity to drop size is an artefact, not a principle to be copied